

1 hex digit is 4 bits.

$$16 = 2^4$$

System A block size 16 byte. 4 bits.

16 elements, each 1 byte.

① 0x3214 → 0x321f
0x3210

② 0x31e8 → 0x31ef
0x31e0

③ 0x322f
0x322f

$$\Rightarrow \frac{4}{10} = 0.4$$

B 256 byte. $\Rightarrow 2^8$ 8 bits

256 elements

\downarrow
 $0x3214$ $0x31e8$ $0x31e8$ $0x3220$ $0x31dc$ $0x31e0$ $0x320c$ $0x31de$ $0x3224$ $0x31f8$
 miss- miss hit hit hit hit hit hit hit hit

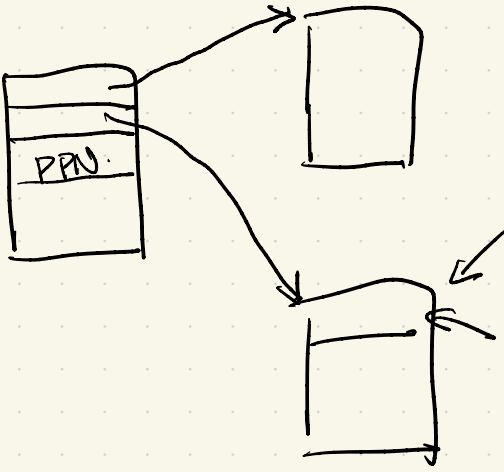
$0x32ff$ } 256 bytes.
 $0x3200$ }
 $\Delta\Delta$

$32 \Rightarrow 2^5$

hit ratio = $\frac{8}{10} = 0.8$

$0x3214 \Rightarrow 0x32$ $\overbrace{0001}$ $\overbrace{0100}$
 Δ

$0x321f$ } 32 bytes.
 $0x3200$ }



2 level page table.

PTE.

$$\frac{\text{page table size}}{\text{PTE}} = \# \text{ entries.}$$

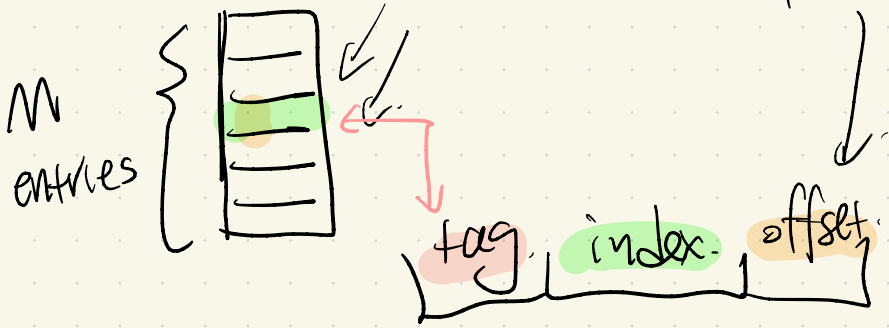
$$\log_2(\# \text{ entries}) = \text{len of VPN}_i.$$

$\{ \text{VPN}_0, \text{VPN}_1, \text{VPN}_2, \dots \}$ offset

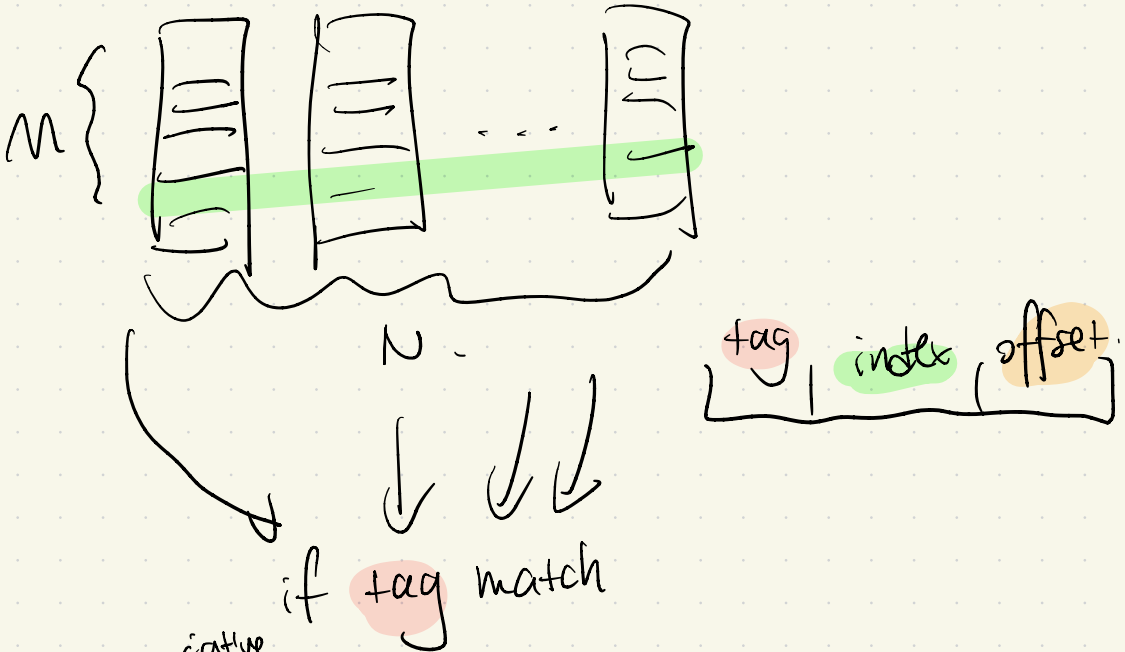


$$\frac{\text{page table size}}{\text{len(PPN)}}$$

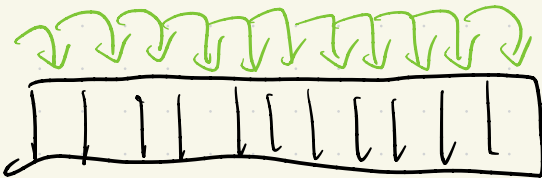
direct-map

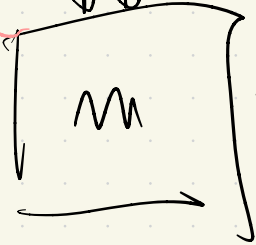
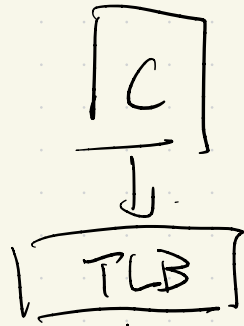
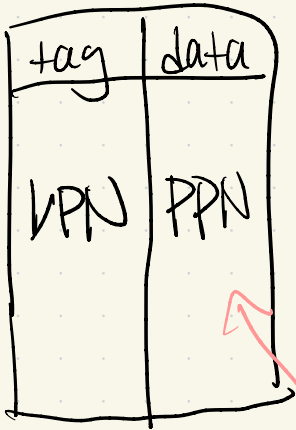


N-way associative

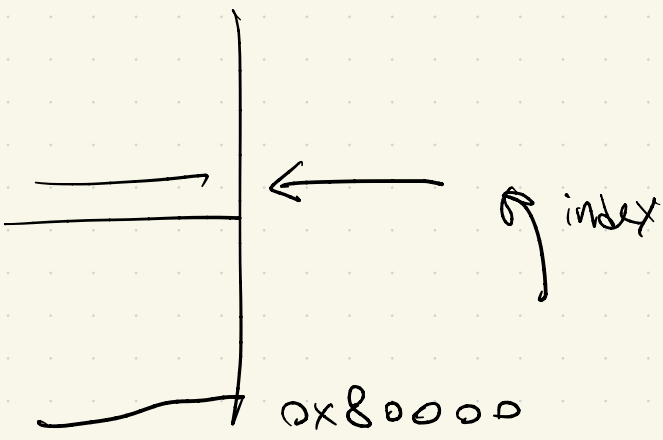
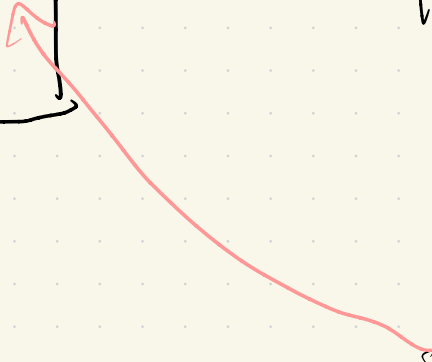


fully associative





times of your page table level



- The page size is 64KiB
- The physical address size is 32 bits
- The virtual address size is 26 bits
- There is a two level page table (each level is the same width)
- All PTEs are 32 bits.
- The internal PTEs only contains the address for the next level
- The leaf PTEs only contain the PPN, not the full address

→ SATP = 0x84816c00

→ effective addr.

0x00f60590
VPN offset

Physical address 4 bytes of data

0xfe1f87d4	0x94d7edf8
0xfe1f3c70	0x00002e42
0xa1a6f854	0x0000201e
0xa1a6dd20	0xd8d0c54c
0xa06b9904	0xdc72cbd0
0xa06b4868	0x00002558
0x84816c24	0xa1a6f800
③ 0x84816c1c	→ 0x77cc3c00 <small>PTE</small>
0x84816c18	0xfe1f3c00
0x84816c08	0x54728c00
0x84816c04	0xa06b4800
④ 0x77cc3c58	→ 0x0000a98 <small>PPN</small>
0x77cbf328	0xeb581204
0x64e10590	0xbec813f8
0x5472f62c	0x6da263ec
0x54728c40	0x000039f8
0x44250590	0x805225b4
0x39f80590	0xc7727140
0x2e420590	0xe5af4aa8
0x25580590	0x108e67d4
0x201e0590	0x31094cdc
0x13320590	0xb05fda58
① 0x0a980590	→ 0x04950a20
0x01980590	0x81d2a7e4

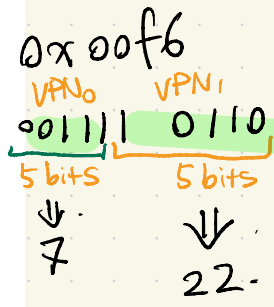
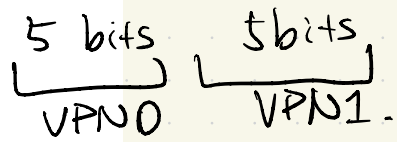
① calculate the offset for page.

$2^{16} \Rightarrow 16 \text{ bits}$

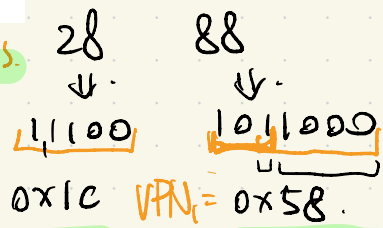
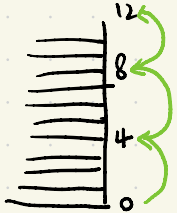
② find the VPN

$\text{len}(\text{VAddr}) = 26$

$\text{len}(\text{index}) = 26 - 16 = 10 \text{ bits}$



★ 4 bytes stored starts that address.



③ find the first page table entry.

$$\text{SATP} = 0x84816c00$$

$$\text{VPN}_0 = 0x1c \Rightarrow \underline{0x84816c1c}$$

④ find the physical page number.

$$\text{PTE} = 0x77cc3c00$$

$$\text{VPN}_1 = 0x58$$

$$\Rightarrow \underline{0x77cc3258}$$

⑤ find Physical address.

$$\text{PPN} = 0xa98$$

$$\text{offset} = 0x0590$$

$$\Rightarrow \text{Paddr} \\ \underline{0xa980590}$$

$$\text{⑥ answer} = \underline{0x04950a20}$$

Q4.

$$t = I \times \text{CPI} \times \frac{1}{f}$$

$$t_1 = I \times \text{CPI}_1 \times \frac{1}{f_1}$$

1 cone

$$t_2 = I \times \text{CPI}_2 \times \frac{1}{f_2}$$

2 cones

$$\text{CPI}_2 = \text{CPI}_1 \times \left(1 - P + \frac{P}{c}\right)$$

$$\frac{1}{f_2} = \frac{1}{f_1} \times \frac{f_1}{f_2}$$

$$I \times \text{CPI}_2 \times \frac{1}{f_2} = I \times \text{CPI}_1 \times \left(1 - P + \frac{P}{c}\right) \times \frac{1}{f_1} \times \frac{f_1}{f_2}$$

$$t_2 = t_1 \times \left(1 - P + \frac{P}{c}\right) \times \frac{f_1}{f_2}$$



$$AMAT = \underbrace{HT} + \underbrace{MP.}$$

L1 \rightarrow 2 cycles.

80% hit ratio.

mem \rightarrow 50 cycles.

$$\underbrace{80\% \times 2 \text{ cycle.}} + \underbrace{20\% (2 + 50)}$$

$$= 80\% \times 2 + 20\% \times 2 + 20\% \times 50.$$

$$= 2 + 20\% \times 50.$$